Amendments to the Claims

1	1. (Currently Amended) A method of performing a transaction in a databa	ise
2	system, comprising:	
3	receiving a transaction to be performed, wherein the transaction is	
4	processed by a plurality of access modules; and	
5	performing a flush of a transaction log from volatile storage to non-	
6	volatile storage by each access module before execution of an end transaction procedu	ıre
7	begins.	
1	2. (Previously Presented) The method of claim 1, further comprising issue	ing
2	a request to flush the transaction log with a message sent to each access module for	
3	performing a last step of the transaction, the last step performed prior to the end	
4	transaction procedure.	
1	3. (Previously Presented) The method of claim 2, further comprising	
2	performing the flush of the transaction log in a data access step prior to the end	
3	transaction procedure to avoid performance of a transaction log flush in the end	
4	transaction procedure.	
1	4. (Previously Presented) The method of claim 2, further comprising	
2	determining that the last step is being performed by all of the plurality of access mode	ales
3	involved in the transaction.	
1	5. (Original) The method of claim 1, further comprising determining if the	ıe
2	transaction log has been flushed before performing the end transaction procedure.	
1	6. (Original) The method of claim 5, further comprising avoiding	
2	performance of a transaction log flush in the end transaction procedure if the transact	ion
3	log has been flushed.	

1	7.	(Original) The method of claim 1, further comprising:
2		identifying the transaction as an implicit transaction.
1	8.	(Original) The method of claim 1, further comprising:
2		performing the end transaction procedure, which follows execution of the
3	transaction.	
	0	(Original) The method of claim 8 merferming the and transportion
1	9.	(Original) The method of claim 8, performing the end transaction
2	procedure co	
3		skipping broadcast of a directive indicating commencement of the end
4	transaction p	rocedure to the plurality of access modules.
1	10.	(Original) A method of performing an end transaction procedure in a
2	database syst	em, comprising:
3		a first access module in the database system writing an end transaction
4	indication to	a first transaction log portion, the first access module being part of a cluster
5	of access mo	dules; and
6		the first access module sending an end transaction directive to a fallback
7	module asso	ciated with the first access module, the fallback module being part of the
8	cluster.	
1	11.	(Original) The method of claim 10, wherein the first access module sends
2		action directive to the fallback module but not to other access modules in the
3	cluster.	
1	12.	(Original) The method of claim 10, wherein sending the end transaction
2	directive con	nprises sending an end transaction-part one directive.
1	13.	(Original) The method of claim 12, further comprising the first access
2	module broa	deasting an end transaction-part two directive to all access modules in the
3	cluster.	

1	14.	(Original) The method of claim 10, further comprising the fallback		
2	module writi	ng an end transaction indication to a second transaction log portion.		
1	15.	(Previously Presented) The method of claim 10, further comprising the		
2	first access n	nodule flushing the first transaction log portion from volatile storage to no		
3	volatile stora	nge.		
1	16.	(Original) The method of claim 10, further comprising the first access		
2	module flush	ning the first transaction log portions but the other access modules in the		
3	cluster not flushing their respective transaction log portions.			
1	17.	(Currently Amended) A database system comprising:		
2		a plurality of storage media, the storage media comprising persistent		
3	storage;			
4		volatile storage; and		
5		a plurality of access modules, wherein each access module is coupled to		
6	one of the pl	urality of storage media; and		
7		each of the access modules being adapted to flush a transaction log from		
8	the volatile	the volatile storage to the persistent storage before performing execution of an end		
9	transaction p	procedure <u>begins</u> .		
1	18.	(Original) The database system of claim 17, further comprising a		
2	controller ad	lapted to determine if each access module has flushed the transaction log		
3	maintained l	by the access module.		
1	19.	(Original) The database system of claim 18, wherein the controller is		
2	adapted to sl	kip sending a directive to perform a transaction log flush if the controller		

determines that each access module has flushed the transaction log before the end

transaction procedure.

	1	20. (Previously Presented) The database system of claim 17, further	
	2	comprising a controller adapted to provide a flush directive with a message to each of the	
	3	access modules to perform a last step of the transaction before the end transaction	
	4	procedure.	
	1	21. (Currently Amended) An article comprising a medium storing instructions	
	2	for enabling a processor-based system to:	
•	3	receive a transaction to be performed, wherein the transaction is processed	
	4	by a plurality of access modules;	
	5	determine that a last step of the transaction involves the plurality of access	
	6	modules, wherein the last step is performed before execution of an end transaction	
	7	procedure begins; and	
	8	flush a transaction log from volatile storage to a non-volatile storage while	
	9	the last step is performed by the plurality of access modules.	
	1	22. (Previously Presented) The article of claim 21, further storing instructions	
	2	for enabling the processor-based system to:	
	3	perform the end transaction procedure, wherein the end transaction	
	4	procedure follows execution of the last step of the transaction.	
	1	23. (Previously Presented) The article of claim 22, further storing instructions	
	2	for enabling a processor-based system to:	
	3	avoid broadcast of a directive indicating commencement of the end	
	4	transaction procedure to the plurality of access modules.	
	1	24. (Previously Presented) A method of performing a transaction in a database	
	2	system, comprising:	
	3	receiving a transaction to be performed on plural access modules in the	
annote the second plantage and the second	4	database system;	
	5	maintaining a log in volatile storage to track operations performed in the	
	6	transaction; and	

7		writing the log to persistent storage before start of an end transaction
8	procedure.	
1	25.	(Original) The method of claim 24, wherein writing the log to persistent
2	storage comp	orises flushing the log.
	26	(O in the D. The country of the Land and the Company of the Land and the Company of the Company
1	26.	(Original) The method of claim 24, wherein maintaining the log comprises
2	maintaining	a transaction log.
1	27.	(Original) The method of claim 24, further comprising performing the end
2	transaction p	rocedure, the end transaction procedure comprising writing an end
3	transaction is	ndication into the log.
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1	28.	(Currently Amended) A database system comprising:
2		storage media comprising persistent storage;
3		volatile storage;
4		access modules coupled to the storage media; and
5		a parsing engine coupled to the access modules, the parsing engine
6	adapted to p	erform one of:
7		(a) providing a directive with a message to perform a last step
8	of a transact	ion and communicating the directive to the access modules, each access
9	module resp	onsive to the directive to perform a transaction log flush from the volatile
10	storage to th	e persistent storage before performance execution of an end transaction
11	procedure be	egins; and
12		(b) determining if each of the access modules has performed a
13	transaction 1	og flush before start of the end transaction procedure;
14		the parsing engine adapted to avoid sending a broadcast directive to the
15	access modu	ales to cause performance of a transaction log flush during the end transaction
16	procedure.	

1	29.	(Previously Presented) The method of claim 1, wherein the transaction
2	comprises plu	aral steps, the method further comprising:
3	•	performing the plural steps prior to performing the end transaction
4	procedure, an	d
5		wherein performing the flush of the transaction log comprises performing
6	the flush of th	ne transaction log in one of the plural steps.
1	30.	(Previously Presented) The method of claim 29, wherein performing the
2		omprises performing, in each of the plural steps, access of relational table
3	data stored in the database system.	
1	31.	(Previously Presented) The method of claim 30, wherein performing the
2	flush of the tr	ransaction log in one of the plural steps comprises performing the flush of
3	the transactio	n log in a last one of the plural steps.
1	32.	(Previously Presented) The method of claim 31, further comprising each
2	access modul	e adding a first entry to the transaction log to redo the transaction by the
3	access modul	e in case of system failure.
1	33.	(Previously Presented) The method of claim 4, wherein performing the
2	flush of the tr	ansaction is prior to the end transaction procedure if the last step is
3	performed by	all of the plurality of access modules, the method further comprising:
4		performing the flush of the transaction log in the end transaction
5	procedure if t	he last step is not performed by all of the plurality of access modules.
	2.4	
1	34.	(Previously Presented) The database system of claim 17, wherein the
2	•	es to perform a transaction comprising plural steps, one or more of the
3		es adapted to perform the plural steps prior to the end transaction procedure
4	and the acces	s modules adapted to perform the flush of the transaction log in one of the
5	plural steps.	

- 1 35. (Previously Presented) The database system of claim 34, wherein the one of the plural steps comprises a last one of the steps.
- 1 36. (Previously Presented) The database system of claim 35, wherein the 2 transaction log comprises a first entry associated with each access module to enable a 3 redo of the transaction in case of system failure.
- 1 37. (Previously Presented) The database system of claim 36, wherein the 2 transaction log further comprises a second entry associated with each access module to 3 enable an undo of the transaction.

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and

- 38. (Previously Presented) The database system of claim 34, further comprising a controller to determine whether a last one of the steps involves all the access modules, and in response to determining that the last one of the steps involves all the access modules, the controller to send a directive to all the access modules to perform the flush of the transaction log in the last one of the steps.
- 39. (Previously Presented) The database system of claim 38, in response to determining that the last step does not involve all access modules, the controller to send a directive to perform the flush of the transaction log in the end transaction procedure.
- 40. (Previously Presented) The article of claim 21, wherein the transaction comprises plural steps, the article further storing instructions for enabling a processor-based system to:
- 4 perform the plural steps prior to performing the end transaction procedure,
- wherein performing the flush of the transaction log comprises performing
 the flush of the transaction log in one of the plural steps.

41. (Previously Presented) The article of claim 40, wherein performing the 2 plural steps comprises performing, in each of the plural steps, access of relational table data stored in the database system. 3

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- 42. (Previously Presented) The article of claim 41, wherein performing the 1 flush of the transaction log in one of the plural steps comprises performing the flush of 2 the transaction log in a last one of the plural steps. 3
- 43. (Previously Presented) The article of claim 42, further storing instructions 1 for enabling a processor-based system to cause each access module to add a first entry to 2 the transaction log to redo the transaction by the access module in case of system failure. 3